



Computational studies of entanglement and quantum contextuality properties towards their formal verification

Henri de Boutray

Femto-ST laboratory

DISC department

VESONTIO team

ANR project: I-QUINS

`henri.de_boutray@univ-fcomte.fr`

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Studies of quantum entanglement and contextuality



Problematic:

- ▶ Lack of specification and verification in current quantum computing

Intermediary objective:

- ▶ Understanding key properties to be specified

PhD Objectives:

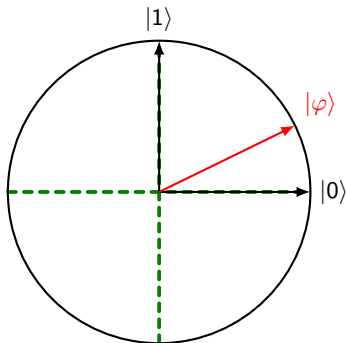
- ▶ Entanglement detection
- ▶ Algorithm specification with entanglement
- ▶ Contextuality detection
 - ▶ Contextual experiment generation



- 1 Background
 - Quantum computing
 - Classical verification
- 2 Entanglement
 - Basics
 - Entanglement evaluation
 - Grover's algorithm
 - Results
- 3 Contextuality
 - Basics
 - Symplectic space
 - A linear problem
 - Results
- 4 Conclusion



The qubit, superposition



$$|0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

$$|1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$$\begin{cases} |\varphi\rangle = \alpha |0\rangle + \beta |1\rangle \\ (\alpha, \beta) \in \mathbb{C}^2, |\alpha|^2 + |\beta|^2 = 1 \end{cases}$$



Qubit transformation: the unitary

$$|\varphi\rangle \longrightarrow \boxed{M} \longrightarrow |\varphi'\rangle$$

$$|\varphi'\rangle = \llbracket M \rrbracket |\varphi\rangle$$

$$M \text{ unitary: } M\bar{M}^T = I$$



Qubit transformation: the unitary

$$|\varphi\rangle \text{ --- } \boxed{M} \text{ --- } |\varphi'\rangle$$

$$M \text{ unitary: } M\bar{M}^T = I$$

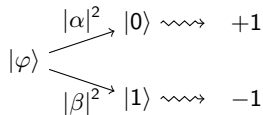
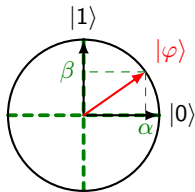
$$|\varphi'\rangle = \llbracket M \rrbracket |\varphi\rangle$$

$$|\varphi\rangle \text{ --- } \boxed{M_1} \text{ --- } \boxed{M_2} \text{ --- } |\varphi'\rangle$$

$$|\varphi'\rangle = M_2 M_1 |\varphi\rangle$$

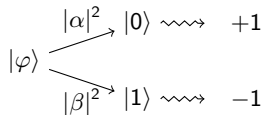
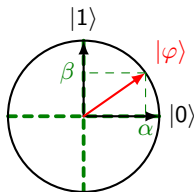


Qubit transformation: the measure





Qubit transformation: the measure



Pauli matrices:

X	Y	Z
$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$	$\begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}$	$\begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$
$\begin{matrix} -\rangle & +\rangle \\ -1 & 1 \end{matrix}$	$\begin{matrix} y_-\rangle & y_+\rangle \\ -1 & 1 \end{matrix}$	$\begin{matrix} 1\rangle & 0\rangle \\ -1 & 1 \end{matrix}$

$$|\pm\rangle = \frac{|0\rangle \pm |1\rangle}{\sqrt{2}}, \quad |y_{\pm}\rangle = \frac{|0\rangle \pm i|1\rangle}{\sqrt{2}}$$

$$E(M_O(|\varphi\rangle)) = \langle \varphi | O | \varphi \rangle$$



Combining qubits

$$|\varphi_1\rangle = \alpha_1 |0\rangle + \beta_1 |1\rangle, \quad |\varphi_2\rangle = \alpha_2 |0\rangle + \beta_2 |1\rangle$$

$$\left. \begin{array}{l} |\varphi_1\rangle \text{ ---} \\ |\varphi_2\rangle \text{ ---} \end{array} \right\} |\varphi_1\rangle \otimes |\varphi_2\rangle$$

$$\begin{aligned} |\varphi_1\rangle \otimes |\varphi_2\rangle &= (\alpha_1 |0\rangle + \beta_1 |1\rangle) \otimes (\alpha_2 |0\rangle + \beta_2 |1\rangle) \\ &= \alpha_1 \alpha_2 |00\rangle + \alpha_1 \beta_2 |01\rangle + \beta_1 \alpha_2 |10\rangle + \beta_1 \beta_2 |11\rangle \end{aligned}$$

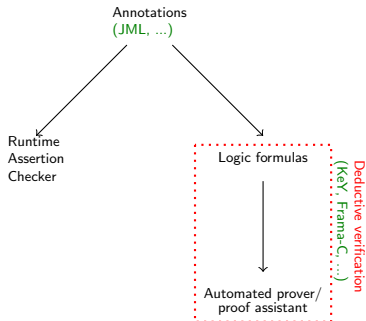
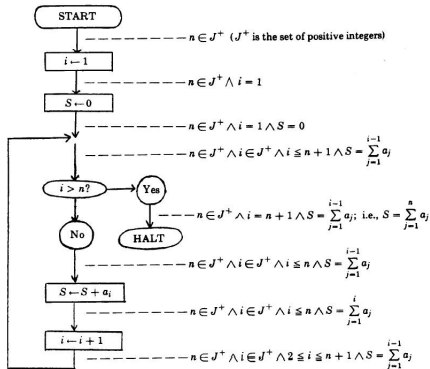
A state on n qubits is a vector of 2^n entries!

Combining gates



$$= (I_4 \otimes M_3 \otimes I_2) \times (M_1 \otimes M_2)$$

State properties



[Flo67] R. Floyd.

Assigning Meanings to Programs.

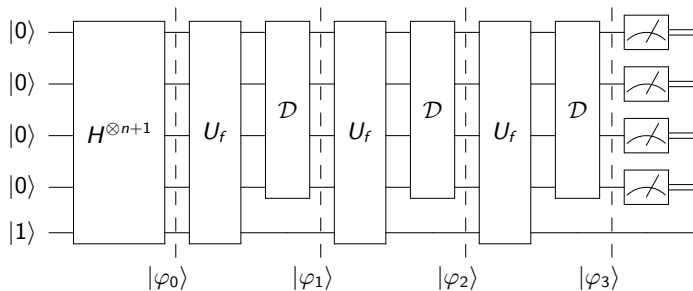
Proceedings of Symposium on Applied Mathematics, 19: 19–32, 1967.



Entanglement

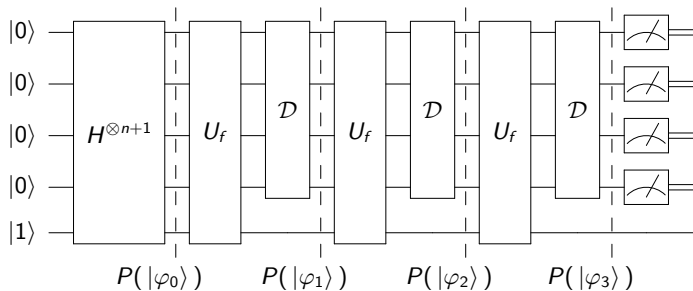


Quantum annotations





Quantum annotations





Bell inequalities (CHSH actually)

Classical :

A	B	A'	B'	$AB + AB' + A'B - A'B'$
-1	-1	-1	-1	2
-1	-1	-1	1	2
-1	-1	1	-1	2
-1	-1	1	1	-2
-1	1	-1	-1	-2
-1	1	-1	1	-2
-1	1	1	-1	2
-1	1	1	1	-2
1	-1	-1	-1	-2
1	-1	-1	1	2
1	-1	1	-1	-2
1	-1	1	1	-2
1	1	-1	-1	-2
1	1	-1	1	2
1	1	1	-1	2
1	1	1	1	2

$$\langle AB \rangle + \langle AB' \rangle + \langle A'B \rangle - \langle A'B' \rangle \leq 2$$

Quantum :

$$|\Phi^-\rangle = \frac{|01\rangle - |10\rangle}{\sqrt{2}}$$

$$A = X, \quad A' = Z, \quad B = -\frac{Z+X}{\sqrt{2}}, \quad B' = \frac{Z-X}{\sqrt{2}}$$

$$\begin{aligned} \langle AB \rangle_{\Phi^-} &= \langle \Phi^- | -X \otimes \frac{Z+X}{\sqrt{2}} | \Phi^- \rangle \\ &= -\frac{1}{2} \langle \Phi^- | (|1\rangle \otimes (|0\rangle - |1\rangle) - \\ &\quad |0\rangle \otimes (|0\rangle + |1\rangle)) \\ &= \frac{1}{\sqrt{2}} \end{aligned}$$

$$\text{similarly } \langle AB' \rangle_{\Phi^-} = \langle A'B \rangle_{\Phi^-} = \frac{1}{\sqrt{2}} \text{ and}$$

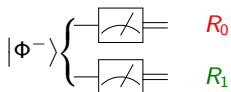
$$\langle A'B' \rangle_{\Phi^-} = -\frac{1}{\sqrt{2}}$$

$$\langle AB \rangle_{\Phi^-} + \langle AB' \rangle_{\Phi^-} + \langle A'B \rangle_{\Phi^-} - \langle A'B' \rangle_{\Phi^-} = 2\sqrt{2} > 2$$

Correlation



$$|\Phi^-\rangle = \frac{|01\rangle - |10\rangle}{\sqrt{2}}$$



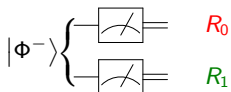
$$R_0 = 1 \implies R_1 = -1$$

$$R_0 = -1 \implies R_1 = 1$$

Correlation



$$|\Phi^-\rangle = \frac{|01\rangle - |10\rangle}{\sqrt{2}}$$



$$R_0 = 1 \implies R_1 = -1$$

$$R_0 = -1 \implies R_1 = 1$$

$|\Phi^-\rangle$ is not separable: $\nexists |\varphi_1\rangle, |\varphi_2\rangle / |\Phi^-\rangle = |\varphi_1\rangle \otimes |\varphi_2\rangle$
 entangled = not separable



Entanglement evaluations

- ▶ entanglement quantification: Geometric Measure of entanglement [WG03], Bell-Mermin inequalities [Mer90, ACG⁺16], Cayley hyperdeterminant [LT03]
- ▶ entanglement classification: Secant varieties [HJN16]

-
- [WG03] Tzu-Chieh Wei and Paul M. Goldbart.
Geometric measure of entanglement and applications to bipartite and multipartite quantum states.
Physical Review A, 68(4):042307, October 2003.
- [Mer90] N David Mermin.
Extreme quantum entanglement in a superposition of macroscopically distinct states.
Physical Review Letters, 65(15):1838–1840, October 1990.
- [ACG⁺16] Daniel Alsina, Alba Cervera, Dardo Goyeneche, José I. Latorre, and Karol Życzkowski.
Operational approach to Bell inequalities: Applications to qutrits.
Physical Review A, 94(3):032102, September 2016.
- [LT03] Jean-Gabriel Luque, and Jean-Yves Thibon.
The Polynomial Invariants of Four Qubits.
Physical Review A, 67, no. 4: 042303, April 2003.
- [HJN16] Frédéric Holweck, Hamza Jaffali, and Ismaël Nounouh.
Grover's algorithm and the secant varieties.
Quantum Information Processing, 15(11):4391–4413, November 2016.



Mermin polynomials

Definition (Mermin polynomials)

Let $a = (a_n)_{n \geq 1}$ and $a' = (a'_n)_{n \geq 1}$ be two families of observables. The *Mermin polynomial* $M_n(a, a')$ is defined by:

$$\begin{cases} M_1(a, a') = a_1 & \text{and} \\ M_n(a, a') = \frac{1}{2} M_{n-1}(a, a') \otimes (a_n + a'_n) + \frac{1}{2} M_{n-1}(a', a) \otimes (a_n - a'_n) & \text{for } n \geq 2 \end{cases}$$



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Example: For two qubits, $M_2 = \frac{1}{2}(a_1 \otimes a_2 + a_1 \otimes a'_2 + a'_1 \otimes a_2 - a'_1 \otimes a'_2)$

Remark: When $a_1 = X$, $a_2 = -\frac{Z+X}{\sqrt{2}}$, $a'_1 = Z$ and $a'_2 = \frac{Z-X}{\sqrt{2}}$, M_2 is the Bell operator.



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To detect entanglement of a given state, we instantiate those Mermin polynomials M_n with specific values of a_n and a'_n .



Mermin evaluation and classical limit

- ▶ Mermin evaluation: $f_{M_n} : |\varphi\rangle \mapsto \langle \varphi | M_n | \varphi \rangle$
- ▶ $|\varphi\rangle$ classical $\implies f_{M_n}(|\varphi\rangle) \leq 1$
- ▶ Mermin evaluation is an entanglement witness

Mermin operator optimization for Grover's algorithm

- ▶ $|\varphi\rangle$ non-local?

Find an M_n such that $f_{M_n}(|\varphi\rangle) > 1$

- ▶ M_n is a function of $(a_i)_{1 \leq i \leq n}$ and $(a'_i)_{1 \leq i \leq n}$

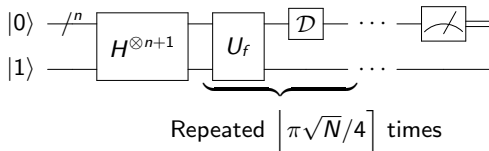
$$\forall i, a_i = \alpha X + \beta Y + \delta Z, \quad a'_i = \alpha' X + \beta' Y + \delta' Z$$

Find $(\alpha, \beta, \delta, \alpha', \beta', \delta')$ such that $f_{M_n}(|\varphi\rangle) > 1$

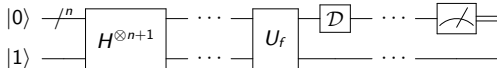
Grover algorithm in a nutshell



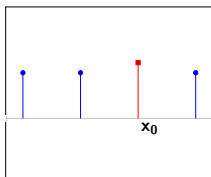
- ▶ Search an item \mathbf{x}_0 in an unsorted database Ω of $N = 2^n$ objects
- ▶ Just by applications of the Boolean function $f : \Omega \rightarrow \{0, 1\}$ such that $f(z) = 1 \Leftrightarrow z = \mathbf{x}_0$
- ▶ $\mathcal{O}(\sqrt{N})$ complexity: quadratic improvement over classical search
- ▶ Oracle U_f defined by $U_f |x, y\rangle = |x, y \oplus f(x)\rangle$
- ▶ Amplitude amplification



Grover's amplitude amplification

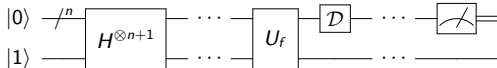


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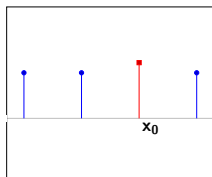


State before U_f

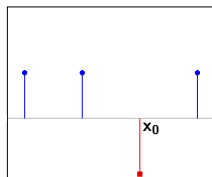
Grover's amplitude amplification



...

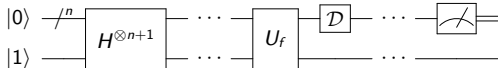


State before U_f

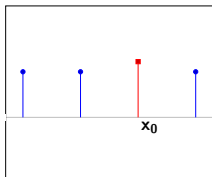


State after U_f

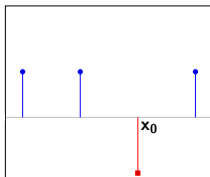
Grover's amplitude amplification



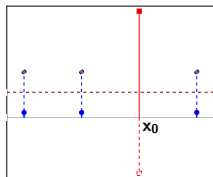
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State before U_f



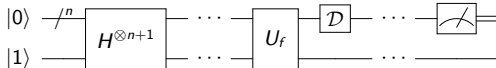
State after U_f



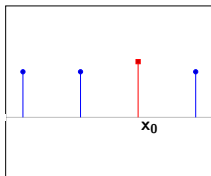
Effect of \mathcal{D}



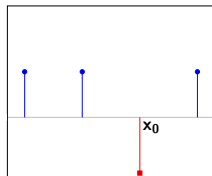
Grover's amplitude amplification



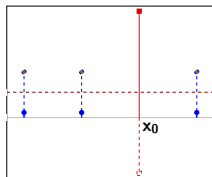
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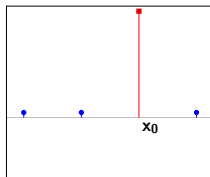
State before U_f



State after U_f



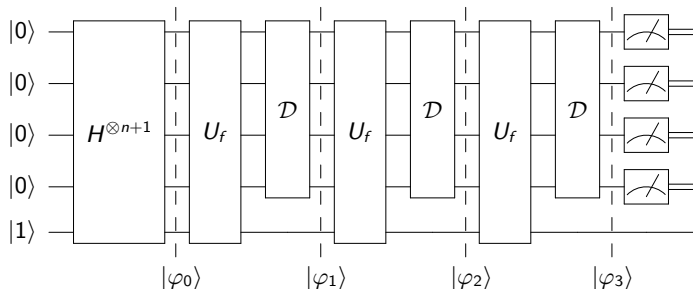
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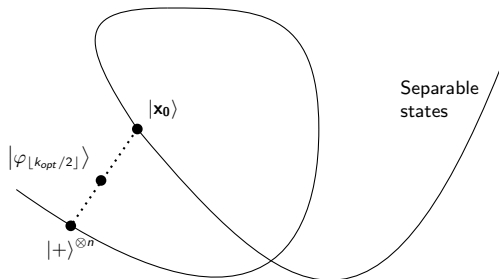
States naming





Preamble

$$|\varphi_k\rangle = \alpha_k |+\rangle^{\otimes n} + \beta_k |\mathbf{x}_0\rangle$$



the middle point is $|\varphi_{ent}\rangle = \frac{|\mathbf{x}_0\rangle + |+\rangle^{\otimes n}}{K}$

$$|\varphi_{[k_{opt}/2]}\rangle \approx |\varphi_{ent}\rangle$$

[HJN16] Frédéric Holweck, Hamza Jaffali, and Ismaël Nounouh.

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Expected entanglement properties

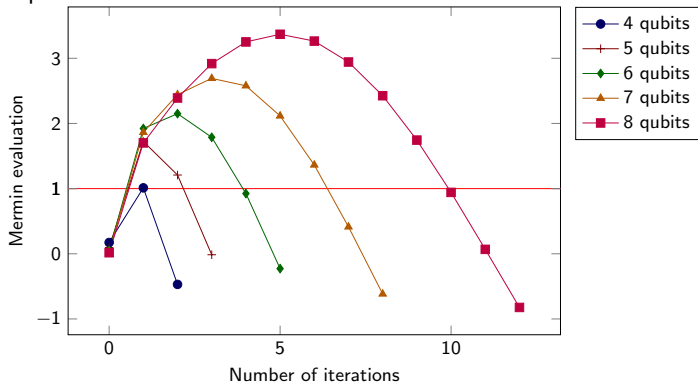
If M_n is chosen to optimize $f_{M_n}(|\varphi_{ent}\rangle)$, then we expect f_{M_n} to behave like a distance measure from $|\varphi_{ent}\rangle$.

Thus we anticipate that:

- ▶ $f_{M_n}(|\varphi_k\rangle)$ reaches maximum around $k_{opt}/2$
- ▶ $f_{M_n}(|\varphi_k\rangle)$ grows for k in $[0, \lfloor k_{opt}/2 \rfloor]$
- ▶ $f_{M_n}(|\varphi_k\rangle)$ decreases for k in $[\lfloor k_{opt}/2 \rfloor + 1, k_{opt}]$

Results, 4 to 8

For 8 qubits, 1 week of computation on personal computer with naive implementation.

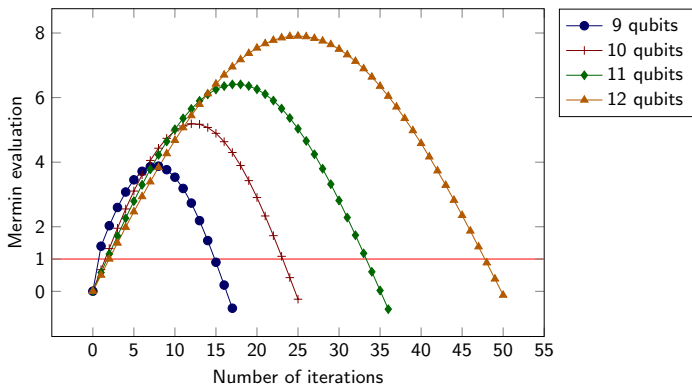


n	4	5	6	7	8
k_{opt}	2	3	5	8	12

Results, 9 to 12



On a supercomputer (Mésocentre UFC):

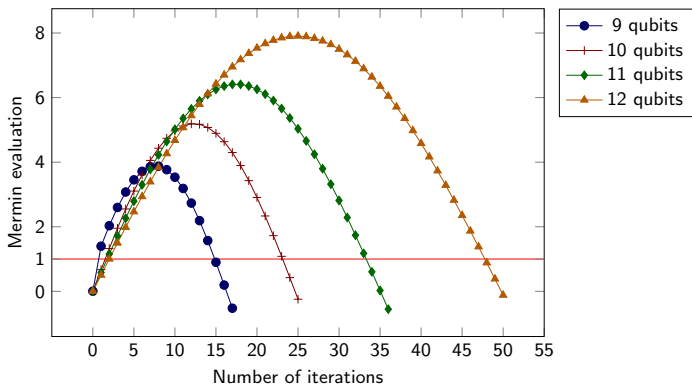


n	9	10	11	12
k_{opt}	17	25	36	50

Results, 9 to 12



On a supercomputer (Mésocentre UFC):



n	9	10	11	12
k_{opt}	17	25	36	50

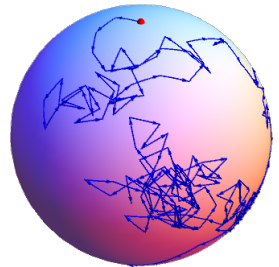
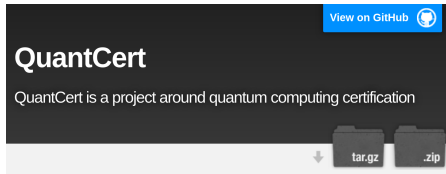
This was in fact a form of RAC!



Technical elements

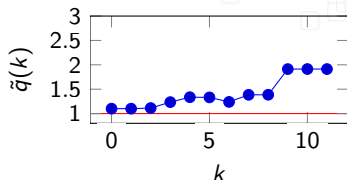


(~ 2000 LoC)

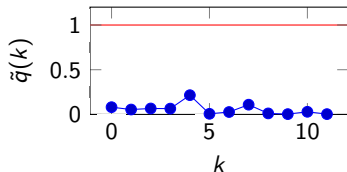


Additional works!

- ▶ Work on the Quantum Fourier Transform



- ▶ Works on IBM's quantum experience (work with Grâce Amouzou, Lomé university, Togo)



- ▶ Formalization of some notions



Coq



QWIRE



Contextuality

The Mermin-Peres square



$$\begin{array}{ccccc}
 X \otimes I & - & I \otimes X & - & X \otimes X \\
 | & & | & & || \\
 I \otimes Y & - & Y \otimes I & - & Y \otimes Y \\
 | & & | & & || \\
 X \otimes Y & - & Y \otimes X & - & Z \otimes Z
 \end{array}$$

The Mermin-Peres square



$$\begin{array}{ccc}
 X \otimes I & - & I \otimes X & - & X \otimes X & / \\
 | & & | & & || & \\
 I \otimes Y & - & Y \otimes I & - & Y \otimes Y & / \\
 | & & | & & || & \\
 X \otimes Y & - & Y \otimes X & - & Z \otimes Z & / \\
 / & & / & & -/ &
 \end{array}$$

The Mermin-Peres square



- 1

$$\begin{array}{ccc}
 X \otimes I & - & I \otimes X & - & X \otimes X & / \\
 | & & | & & || & \\
 I \otimes Y & - & Y \otimes I & - & Y \otimes Y & / \\
 | & & | & & || & \\
 X \otimes Y & - & Y \otimes X & - & Z \otimes Z & / \\
 / & & / & & - / &
 \end{array}$$

The Mermin-Peres square



$$\begin{array}{ccc}
 -1 & & -1 \\
 X \otimes I & - & I \otimes X & - & X \otimes X & / \\
 | & & | & & || & \\
 I \otimes Y & - & Y \otimes I & - & Y \otimes Y & / \\
 | & & | & & || & \\
 X \otimes Y & - & Y \otimes X & - & Z \otimes Z & / \\
 / & & / & & -/ &
 \end{array}$$

The Mermin-Peres square

$$\begin{array}{ccc}
 -1 & -1 & 1 \\
 X \otimes I - I \otimes X - X \otimes X & / \\
 | & | & || \\
 I \otimes Y - Y \otimes I - Y \otimes Y & / \\
 | & | & || \\
 X \otimes Y - Y \otimes X - Z \otimes Z & / \\
 / & / & -/
 \end{array}$$



The Mermin-Peres square



$$\begin{array}{ccc}
 -1 & -1 & 1 \\
 X \otimes I - I \otimes X - X \otimes X & / & \\
 | & | & || \\
 1 & & \\
 I \otimes Y - Y \otimes I - Y \otimes Y & / & \\
 | & | & || \\
 X \otimes Y - Y \otimes X - Z \otimes Z & / & \\
 / & / & -/
 \end{array}$$

The Mermin-Peres square



$$\begin{array}{ccc}
 -1 & -1 & 1 \\
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 | & | & || \\
 1 & 1 & \\
 I \otimes Y - Y \otimes I - Y \otimes Y & / & \\
 | & | & || \\
 X \otimes Y - Y \otimes X - Z \otimes Z & / & \\
 / & / & -/
 \end{array}$$

The Mermin-Peres square



$$\begin{array}{ccc}
 -1 & -1 & 1 \\
 X \otimes I - I \otimes X - X \otimes X & / & \\
 | & | & || \\
 1 & 1 & 1 \\
 I \otimes Y - Y \otimes I - Y \otimes Y & / & \\
 | & | & || \\
 X \otimes Y - Y \otimes X - Z \otimes Z & / & \\
 / & / & -/
 \end{array}$$

The Mermin-Peres square

$$\begin{array}{ccc}
 -1 & -1 & 1 \\
 X \otimes I - I \otimes X - X \otimes X & / & \\
 | & | & || \\
 1 & 1 & 1 \\
 I \otimes Y - Y \otimes I - Y \otimes Y & / & \\
 | & | & || \\
 -1 & & \\
 X \otimes Y - Y \otimes X - Z \otimes Z & / & \\
 / & / & -/
 \end{array}$$



The Mermin-Peres square



$$\begin{array}{ccc}
 -1 & -1 & 1 \\
 X \otimes I - I \otimes X - X \otimes X & / & \\
 | & | & || \\
 1 & 1 & 1 \\
 I \otimes Y - Y \otimes I - Y \otimes Y & / & \\
 | & | & || \\
 -1 & -1 & \\
 X \otimes Y - Y \otimes X - Z \otimes Z & / & \\
 / & / & -/
 \end{array}$$

The Mermin-Peres square



$$\begin{array}{ccc}
 -1 & -1 & 1 \\
 X \otimes I - I \otimes X - X \otimes X & / & \\
 | & | & || \\
 1 & 1 & 1 \\
 I \otimes Y - Y \otimes I - Y \otimes Y & / & \\
 | & | & || \\
 -1 & -1 & ? \\
 X \otimes Y - Y \otimes X - Z \otimes Z & / & \\
 / & / & -/
 \end{array}$$

Unsatisfiable!



Contextuality

An experiment is contextual if no non-contextual classical theory can predict its results.

In quantum physics, a context $c \in C$ is a sequence of compatible (commuting) observables. The measures e_O are eigenvalues of observables $O \in \mathcal{O}$ and the product of measures is the eigenvalue of the product of observables:

$$\prod_{O \in c} e_O = e_{\prod_{O \in c} O}.$$

The experiment (\mathcal{O}, C) is contextual if

$$\nexists f : \mathcal{O} \rightarrow \{-1, 1\} / \forall c \in C, \prod_{O \in c} f(O) = e_{\prod_{O \in c} O}.$$



Binary polar symplectic space

We are interested by geometries such as the Mermin-Peres square, we explore a space containing them, the binary polar symplectic space.

$$O = s \otimes_k O_k \in \mathcal{P}_n \text{ with } O_k \in \{I, X, Y, Z\} \text{ and } s \in \{\pm 1, \pm i\}$$

We build the bijection $\pi_n : \mathcal{P}_n / \{\pm I, \pm iI\} \rightarrow (\mathbb{Z}/2\mathbb{Z})^{2n}$

$$\pi_1(I) = (0, 0)$$

$$\pi_1(X) = (0, 1)$$

$$\pi_1(Y) = (1, 1)$$

$$\pi_1(Z) = (1, 0)$$

$$\pi_n(O_1 \otimes O_2) = \text{cat}(\pi_{n_1}(O_1), \pi_{n_2}(O_2)), \quad n_1 + n_2 = n$$

Example: $\pi_3(X \otimes Y \otimes I) = (0, 1, 1, 1, 0, 0)$



Symplectic product

We have a representation of the *operators* (without their phase), of their *product* (the sum of points), but we lost the *commutation relation*. To recover it, let us check the condition on the coefficients.

$$\pi_1(O) = (z, x), \quad O = sZ^z X^x, s \in \{\pm 1, \pm i\}$$

$$\begin{aligned} OO' &= O'O \\ ss'(Z^z X^x)(Z^{z'} X^{x'}) &= s's(Z^{z'} X^{x'})(Z^z X^x) \\ (Z^z X^x)(Z^{z'} X^{x'}) - (Z^{z'} X^{x'})(Z^z X^x) &= 0 \\ (-1)^{xz'} Z^z Z^{z'} X^x X^{x'} - (-1)^{x'z} Z^{z'} Z^z X^{x'} X^x &= 0 \\ ((-1)^{xz'} - (-1)^{x'z}) Z^{z+z'} X^{x+x'} &= 0 \\ ((-1)^{xz'} - (-1)^{x'z}) &= 0 \\ xz' &= x'z \end{aligned}$$

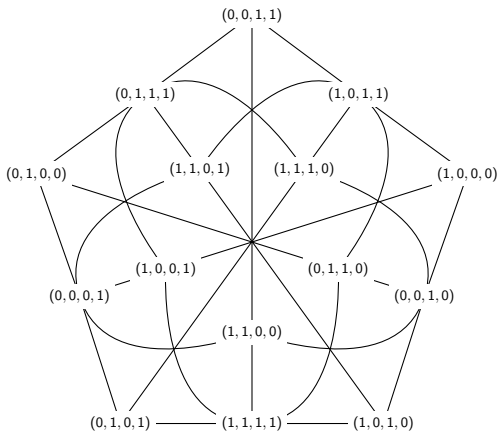
For $\pi_n(O) = (o_1, \dots, o_{2n})$,

$$\langle O|O' \rangle = \sum_{i=1}^n o_{2i-1} o'_{2i} + o_{2i} o'_{2i-1}$$



Building geometries

The space with all possible points and the isotropic subspaces: $W(2n - 1, 2)$ (shorthand W_n). Example for $n = 2$, the Doily:





Hyperplanes

Isotropic subspace of W_n : maximal set of points P such that
 $\forall p_1, p_2 \in P, \langle p_1 | p_2 \rangle = 0$.

The *lines* of W_n are isotropic subspaces of three elements.

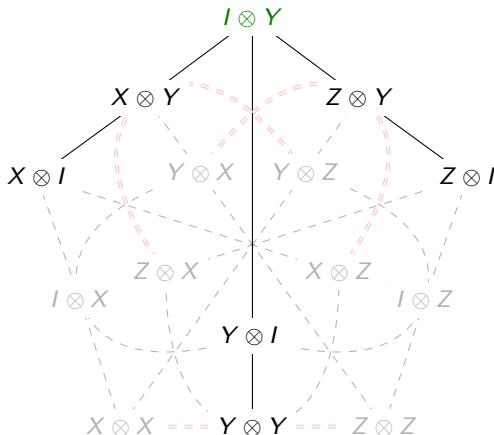
The *hyperplanes* are isotropic subspaces of W_n such that a line of W_n is either entirely in the subspace or intersecting the subspace in a single point.

A hyperplane is either a *perpset* or a *quadric*.

Perpset

A perpset P_r is a set of points that do commute with a single point r :

$$P_r = \{p \in W_n, \langle p|r \rangle = 0\}$$



Quadric



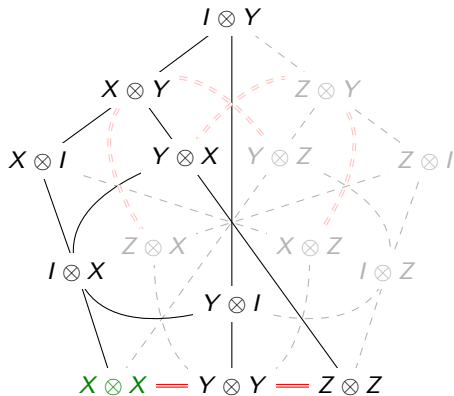
The standard quadratic form $Q_0(x) = \sum_{i=1}^n x_{2i-1}x_{2i}$ let us define a *quadratic form* for each point p of W_n : $Q_p(x) = Q_0(x) + \langle x|p\rangle$.

The *quadric* Q_p is the set of points annihilating the quadratic form Q_p .

A quadric (resp. quadratic form) Q_p (resp. Q_p) is said to be *hyperbolic* if $Q_0(p) = 0$, and *elliptic* otherwise.

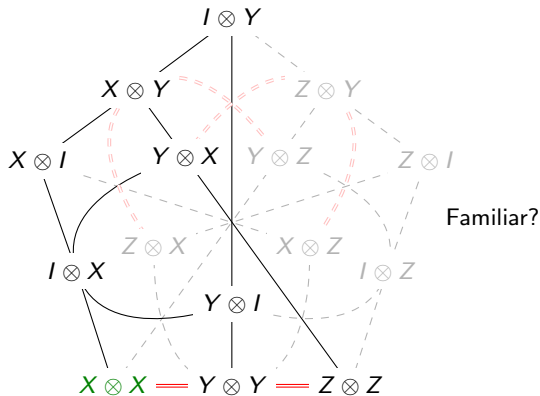
Note that Q_0 is counted in the hyperbolic quadrics.

Hyperbolic



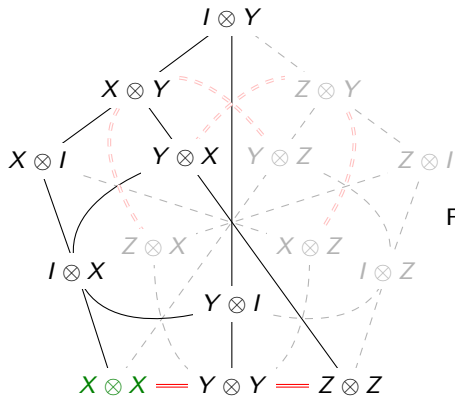


Hyperbolic

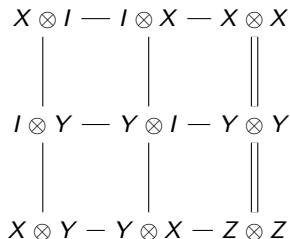




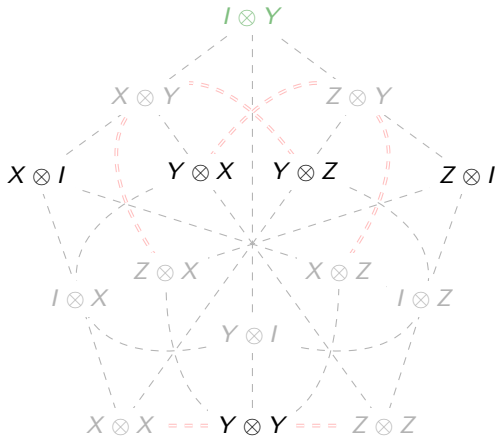
Hyperbolic



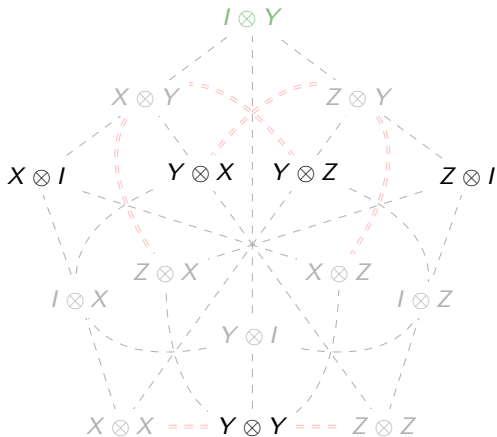
Familiar?



Elliptic



Elliptic



No line! . . . only for W_2



Contextuality as a linear problem

Recall, the experiment $G = (\mathcal{O}, C)$ is contextual when

$$\nexists f : \mathcal{O} \rightarrow \{-1, 1\} / \forall c \in C, \prod_{O \in c} f(O) = e_{\prod_{O \in c} O}.$$

$$(\{-1, 1\}, \times) \rightarrow (\mathbb{Z}/2\mathbb{Z}, +)$$

Let $A = \text{Inc}(G)$, e the evaluation vector such that its c^{th} entry e_c corresponds to context $c \in C$ and $(-1)^{e_c} I = \prod_{O \in c} O$:

$$\nexists x / Ax = e$$

Families contextuality



Geometries	$n = 2$	$n = 3$	$n = 4$	$n = 5$
Whole space	C(1)	C(1)	C(1)	C(1)
Hyperbolics	C(10)	C(36)	C(136)	C(528)
Elliptics	N/A (6)	C(28)	C(120)	C(496)
Perpsets	N(15)	N(63)	N(255)	N(1023)

n : Number of qubits of the system

N: Non-contextual

C: Contextual

(k): There are k instances in this family

N/A: Not applicable

bold: New results



Technicalities

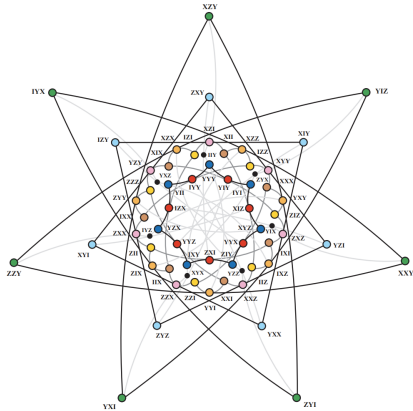
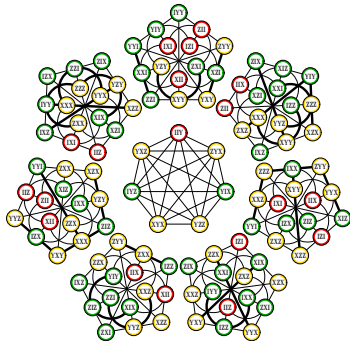
MAGMA
COMPUTER • ALGEBRA

(~ 3000 LoC)

$4^n - 1$ points in $W_n \implies \left\{ \begin{array}{l} \text{Geometric understanding} \\ \text{Intensive computer testing} \end{array} \right.$

- ▶ Preamble work on incidence structures (work lead by Jessy Colonval, UFC),
- ▶ many more geometries explored (with Metod Saniga, Astronomical Institute of the Slovak Academy of Sciences)

Other geometries





Works

- ▶ Journal articles
 - ▶ QIP'20
 - ▶ Mathematics'21
- ▶ Conference paper
 - ▶ AFADL'19
- ▶ Conference posters
 - ▶ GDR-IM'20
 - ▶ QPL'21
- ▶ Presentations
 - ▶ JDD'19
 - ▶ GT-IQ'19
 - ▶ R&Days'21
 - ▶ Auburn algebra seminar'21
- ▶ Website
 - ▶ quantcert.github.io
- ▶ Draft
 - ▶ Contextuality degree study



New horizons

- ▶ Using these notions in formalization frameworks
 - ▶ Mermin polynomials (started in QWIRE, Coq)
 - ▶ Contextuality (started in Why3, Qbricks ?)

- ▶ More geometries to explore!
 - ▶ Cayley hexagons
 - ▶ Quadratic W_4 doilies
 - ▶ ...



Thank you for your attention